

NAME

vm_map - virtual address space portion of virtual memory subsystem

SYNOPSIS

```
#include <sys/param.h>
#include <vm/vm.h>
#include <vm/vm_map.h>
```

DESCRIPTION

The **vm_map** subsystem is used to manage virtual address spaces. This section describes the main data structures used within the code.

The *struct vm_map* is a generic representation of an address space. This address space may belong to a user process or the kernel. The kernel actually uses several maps, which are maintained as subordinate maps, created using the *vm_map_submap(9)* function.

```
struct vm_map {
    struct vm_map_entry header;
    struct sx lock;
    struct mtx system_mtx;
    int nentries;
    vm_size_t size;
    u_int timestamp;
    u_char needs_wakeup;
    u_char system_map;
    vm_flags_t flags;
    vm_map_entry_t root;
    pmap_t pmap;
    int busy;
};
```

The fields of *struct vm_map* are as follows:

<i>header</i>	Head node of a circular, doubly linked list of <i>struct vm_map_entry</i> objects. Each object defines a particular region within this map's address space.
<i>lock</i>	Used to serialize access to the structure.
<i>system_mtx</i>	A mutex which is used if the map is a system map.

<i>nentries</i>	A count of the members in use within the circular map entry list.
<i>size</i>	Specifies the size of the virtual address space.
<i>timestamp</i>	Used to determine if the map has changed since its last access.
<i>needs_wakeup</i>	Indicates if a thread is waiting for an allocation within the map. Used only by system maps.
<i>system_map</i>	Set to TRUE to indicate that map is a system map; otherwise, it belongs to a user process.
<i>flags</i>	Map flags, described below.
<i>root</i>	Root node of a binary search tree used for fast lookup of map entries.
<i>pmap</i>	Pointer to the underlying physical map with which this virtual map is associated.
<i>busy</i>	Map busy counter, prevents forks.

Possible map flags:

MAP_WIREFUTURE	Wire all future pages in this map.
MAP_BUSY_WAKEUP	There are waiters for the map busy status.

The following flags can be passed to `vm_map_find(9)` and `vm_map_insert(9)` to specify the copy-on-write properties of regions within the map:

MAP_COPY_ON_WRITE	The mapping is copy-on-write.
MAP_NOFAULT	The mapping should not generate page faults.
MAP_PREFAULT	The mapping should be prefaulted into physical memory.
MAP_PREFAULT_PARTIAL	The mapping should be partially prefaulted into physical memory.
MAP_DISABLE_SYNCER	Do not periodically flush dirty pages; only flush them when absolutely necessary.

MAP_DISABLE_COREDUMP

Do not include the mapping in a core dump.

MAP_PREFAULT_MADVISE Specify that the request is from a user process calling `madvise(2)`.

MAP_ACC_CHARGED Region is already charged to the requestor by some means.

MAP_ACC_NO_CHARGE Do not charge for allocated region.

The *struct vm_map_entry* is a generic representation of a region. The region managed by each entry is associated with a *union vm_map_object*, described below.

```
struct vm_map_entry {
    struct vm_map_entry *prev;
    struct vm_map_entry *next;
    struct vm_map_entry *left;
    struct vm_map_entry *right;
    vm_offset_t start;
    vm_offset_t end;
    vm_offset_t avail_ssize;
    vm_size_t adj_free;
    vm_size_t max_free;
    union vm_map_object object;
    vm_ooffset_t offset;
    vm_eflags_t eflags;
    /* Only in task maps: */
    vm_prot_t protection;
    vm_prot_t max_protection;
    vm_inherit_t inheritance;
    int wired_count;
    vm_pindex_t lastr;
};
```

The fields of *struct vm_map_entry* are as follows:

prev Pointer to the previous node in a doubly-linked, circular list.

next Pointer to the next node in a doubly-linked, circular list.

left Pointer to the left node in a binary search tree.

<i>right</i>	Pointer to the right node in a binary search tree.
<i>start</i>	Lower address bound of this entry's region.
<i>end</i>	Upper address bound of this entry's region.
<i>avail_ssize</i>	If the entry is for a process stack, specifies how much the entry can grow.
<i>adj_free</i>	The amount of free, unmapped address space adjacent to and immediately following this map entry.
<i>max_free</i>	The maximum amount of contiguous free space in this map entry's subtree.
<i>object</i>	Pointer to the <i>struct vm_map_object</i> with which this entry is associated.
<i>offset</i>	Offset within the <i>object</i> which is mapped from <i>start</i> onwards.
<i>eflags</i>	Flags applied to this entry, described below.

The following five members are only valid for entries forming part of a user process's address space:

<i>protection</i>	Memory protection bits applied to this region.
<i>max_protection</i>	Mask for the memory protection bits which may be actually be applied to this region.
<i>inheritance</i>	Contains flags which specify how this entry should be treated during fork processing.
<i>wired_count</i>	Count of how many times this entry has been wired into physical memory.
<i>lastr</i>	Contains the address of the last read which caused a page fault.

The following flags may be applied to each entry, by specifying them as a mask within the *eflags* member:

MAP_ENTRY_NOSYNC	The system should not flush the data associated with this map periodically, but only when it needs to.
MAP_ENTRY_IS_SUB_MAP	If set, then the <i>object</i> member specifies a subordinate map.
MAP_ENTRY_COW	Indicate that this is a copy-on-write region.

MAP_ENTRY_NEEDS_COPY	Indicate that a copy-on-write region needs to be copied.
MAP_ENTRY_NOFAULT	Specifies that accesses within this region should never cause a page fault. If a page fault occurs within this region, the system will panic.
MAP_ENTRY_USER_WIRED	Indicate that this region was wired on behalf of a user process.
MAP_ENTRY_BEHAV_NORMAL	The system should use the default paging behaviour for this region.
MAP_ENTRY_BEHAV_SEQUENTIAL	The system should depress the priority of pages immediately preceding each page within this region when faulted in.
MAP_ENTRY_BEHAV_RANDOM	Is a hint that pages within this region will be accessed randomly, and that prefetching is likely not advantageous.
MAP_ENTRY_IN_TRANSITION	Indicate that wiring or unwiring of an entry is in progress, and that other kernel threads should not attempt to modify fields in the structure.
MAP_ENTRY_NEEDS_WAKEUP	Indicate that there are kernel threads waiting for this region to become available.
MAP_ENTRY_NOCOREDUMP	The region should not be included in a core dump.

The *inheritance* member has type *vm_inherit_t*. This governs the inheritance behaviour for a map entry during fork processing. The following values are defined for *vm_inherit_t*:

VM_INHERIT_SHARE	The object associated with the entry should be cloned and shared with the new map. A new <i>struct vm_object</i> will be created if necessary.
VM_INHERIT_COPY	The object associated with the entry should be copied to the new map.
VM_INHERIT_NONE	The entry should not be copied to the new map.
VM_INHERIT_DEFAULT	Specifies the default behaviour, VM_INHERIT_COPY.

The *union vm_map_object* is used to specify the structure which a *struct vm_map_entry* is associated

with.

The fields of *union vm_map_object* are as follows:

```
union vm_map_object {
    struct vm_object *vm_object;
    struct vm_map *sub_map;
};
```

Normally, the *sub_map* member is only used by system maps to indicate that a memory range is managed by a subordinate system map. Within a user process map, each *struct vm_map_entry* is backed by a *struct vm_object*.

SEE ALSO

pmap(9), vm_map_check_protection(9), vm_map_delete(9), vm_map_entry_resize_free(9), vm_map_find(9), vm_map_findspace(9), vm_map_inherit(9), vm_map_init(9), vm_map_insert(9), vm_map_lock(9), vm_map_lookup(9), vm_map_madvise(9), vm_map_max(9), vm_map_min(9), vm_map_pmap(9), vm_map_protect(9), vm_map_remove(9), vm_map_simplify_entry(9), vm_map_stack(9), vm_map_submap(9), vm_map_sync(9), vm_map_wire(9)

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